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Hardware attacks and software induced HW attacks and the need for separated trust anchors!

Georg Sigl

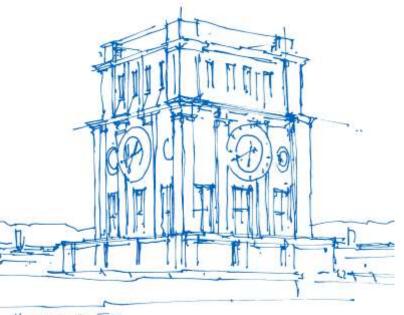
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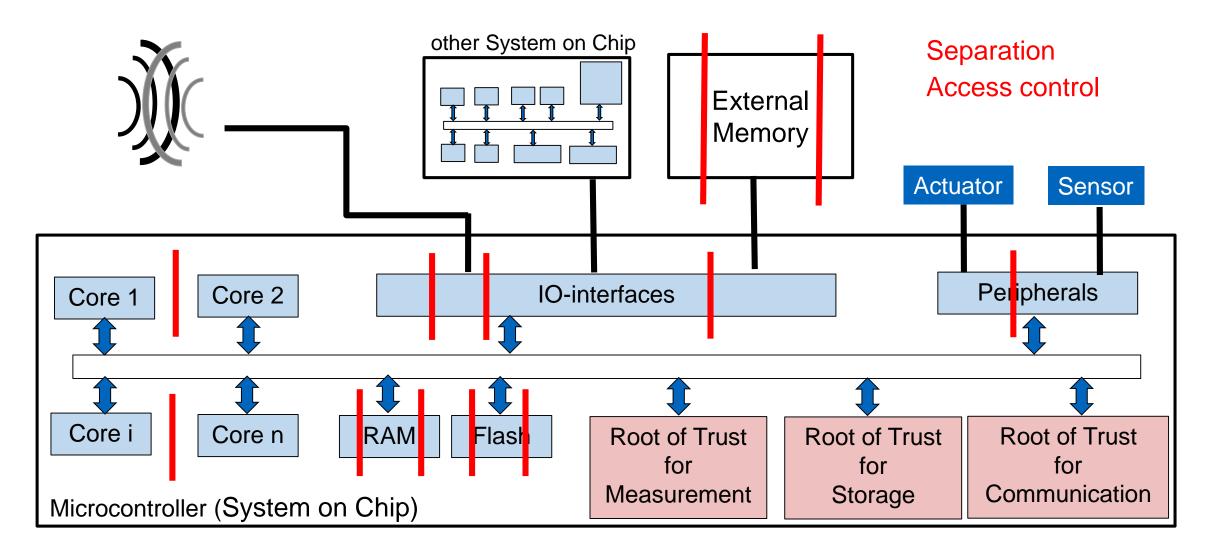
October 2018







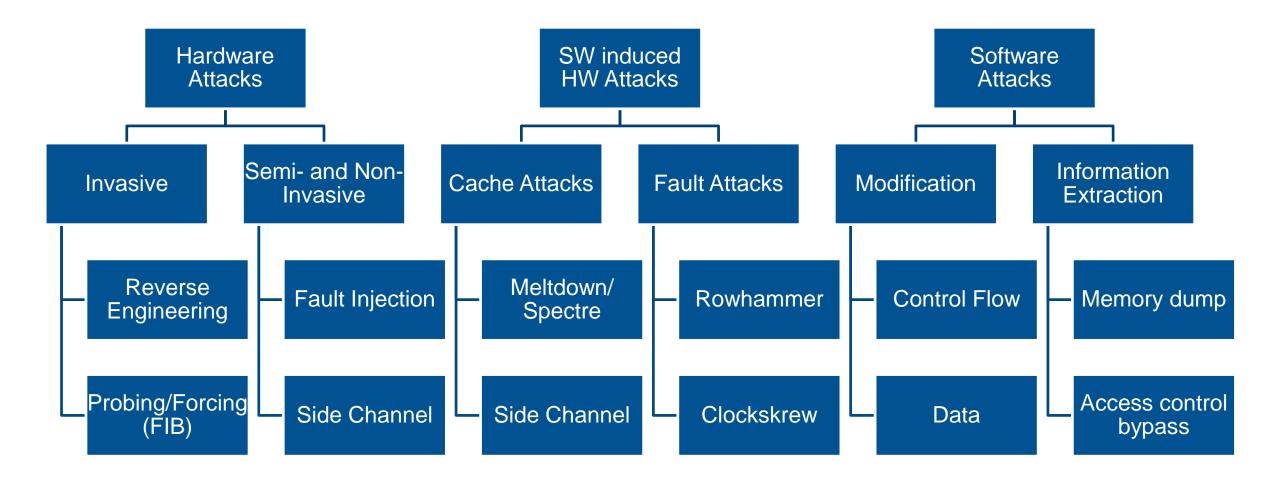
Security in Embedded Systems



Roots of Trust

- Root of Trust for Measurement
 - Program to measure code which runs on system in order to ensure system integrity
 - Program must not be changed by any attack and execution must be in a secure/isolated environment
- Root of Trust for Storage
 - Secret key which can be used to encrypt and authenticate externally stored information
 - Key should never leave SoC
 - Key must not leak to any other party
- Root of Trust for Communication
 - Asymmetric key pair for internet communication
 - Secret part of key pair must not leave SoC
 - Trusted public key of a certification authority

Attack types



High-Resolution Magnetic Field Side-Channels Localized EM against Asymmetric Crypto

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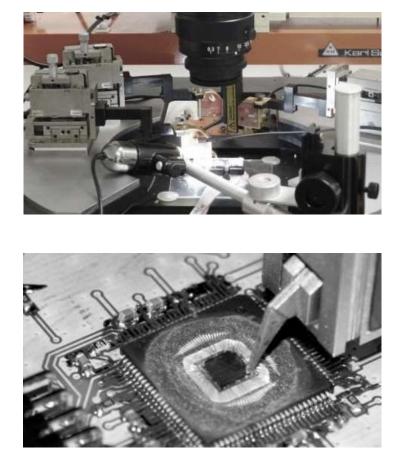
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6:

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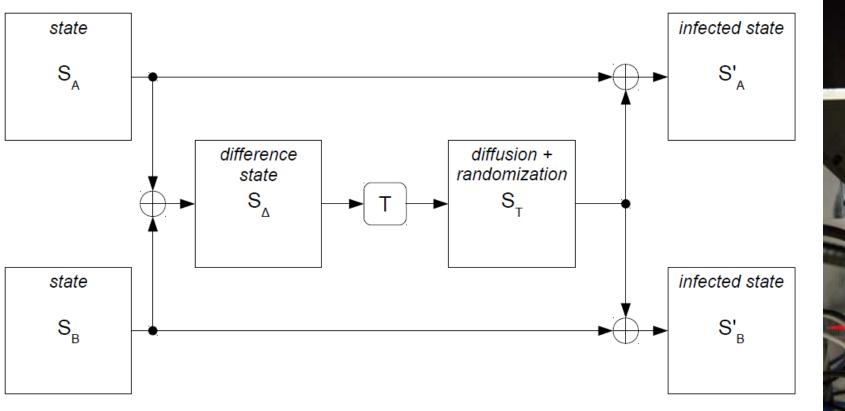


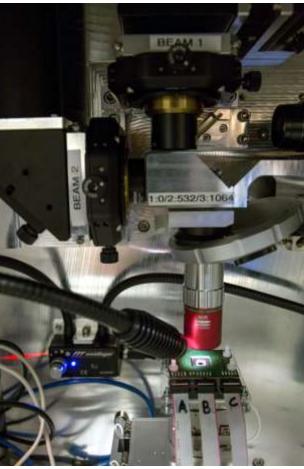
Input: Secret
$$d = d_D d_{D-1} \dots d_2 d_1$$
 with $d_i \in \{0, 1\}$
1: for $i = D$ downto 1 do
2: if $d_i = 1$ then
3: $c \leftarrow c^2 + a$
4: $a \leftarrow c$
5: else
6: $c \leftarrow c^2 + b$
7: $b \leftarrow c$
8: end if
9: end for
integrated circuit $a c$ registers

Heyszl, Johann; Mangard, Stefan; Heinz, Benedikt; Stumpf, Frederic; Sigl, Georg: Localized Electromagnetic Analysis of Cryptographic Implementations. CT-RSA, 2012 Georg Sigl; Chair for Security in Information Technology



Fault Attack on an AES with a simple countermeasure





Selmke, Bodo and Heyszl, Johann and Sigl, Georg: Attack on a DFA protected AES by Simultaneous Laser Fault Injections. FDTC 2016 Fault Diagnosis and Tolerance in Cryptography, 2016

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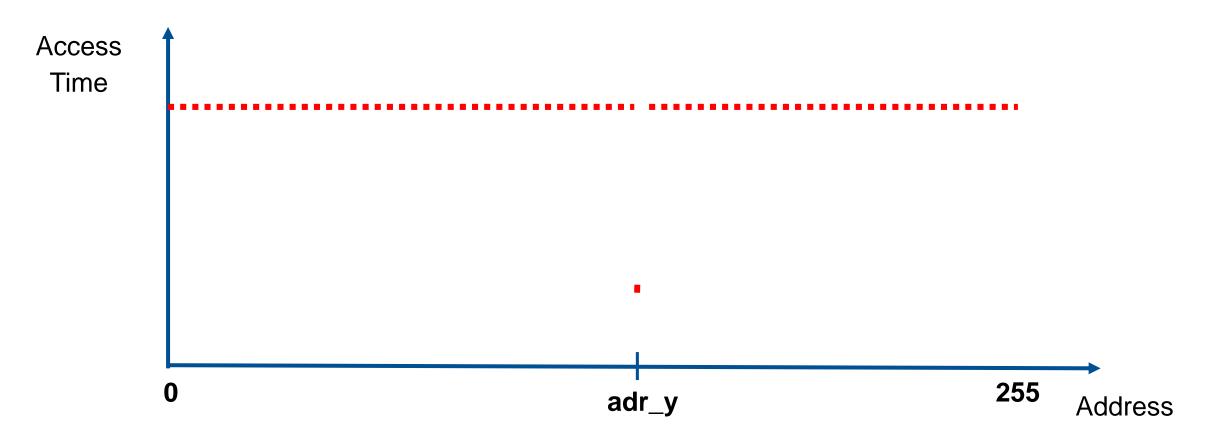
Cache Side Channel Attack: Flush & Reload

Attack Program Attack Program Attack Program **Attack Program Flush** Shared adr_0 Data adr_0 Data data array adr_y Data adr_y Data Data accessed in adr_255 Data cache adr 255 Data Victim Program Victim Program Victim Program Victim Program Access data(adr_y) 2nd access is fast First access takes long

Access data(adr 0 to 255)



Cache Side Channel Attack: Flush & Reload



Cache Attack: Meltdown

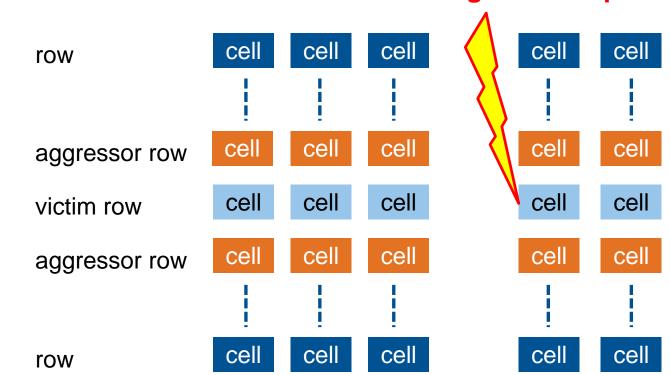
Security: MMU (Memory Management Unit) prevents data access for an attacker to the OS kernel How to circumvent?

- Attacker trains the prefetcher with a dedicated program, e.g. a loop
- Attack program suddenly
 - accesses a byte at a memory address in kernel space
 - \rightarrow CPU prefetcher fetches one byte of data stored at this address in kernel space into cache
 - accesses an **array** with this byte value as index
- Then CPU detects that the branch prediction (with the access right violation) was wrong and reverts most actions, but **cache data are not deleted**
- Now the attacker accesses the complete **array** and measures access time
- The access with the shortest access time corresponds to the read byte value

Root Cause: Wrong assumption during CPU design that speculation can be undone

Software based fault attack: Rowhammer



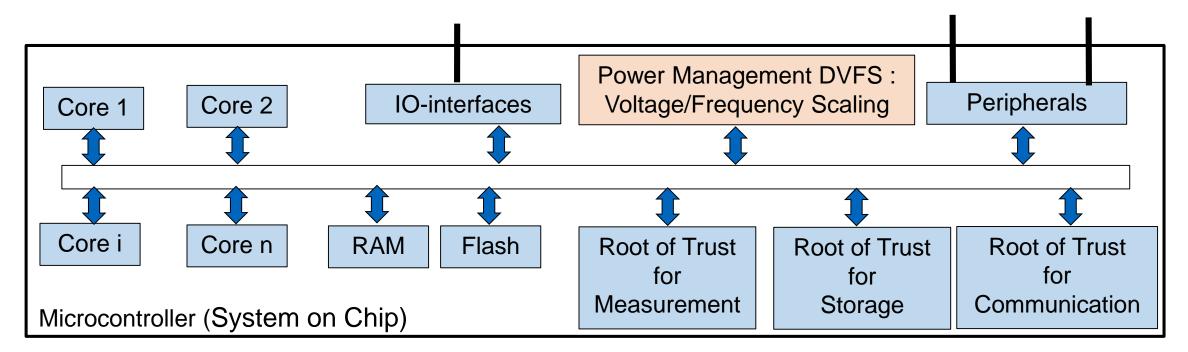


Discharge → bit-flip

- Reading frequently aggressor rows changes contents in victim
- Attacks
 - Change access rights
 - Escape sand boxes (from browser)
 - ...
- Root Cause
- DRAM considers average usage only not an attack
- No error correction implemented (too expensive, performance?)

Flipping bits in memory without accessing them: An experimental study of DRAM disturbance errors. Y. Kim, R. Daly, J. Kim, C.Fallin, J. Lee, D. Lee, C. Wilkerson, K.Lai, O. Mutlu; 2014 ACM/IEEE 41st International Symposium on Computer Architecture (ISCA)

Software based fault attack: CLKskrew

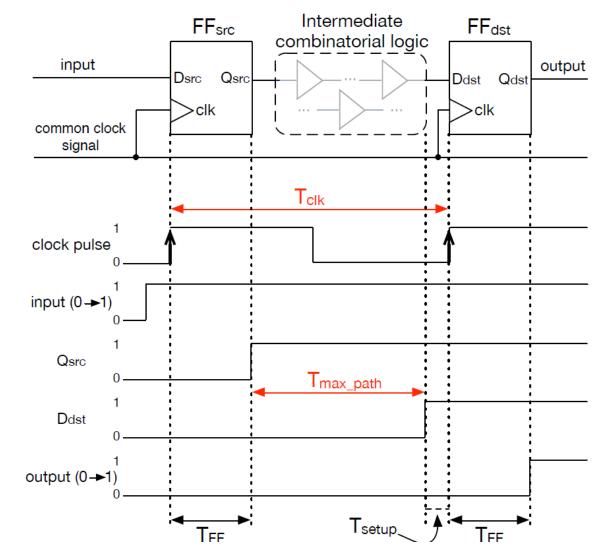


"CLKSCREW: Exposing the Perils of Security-Oblivious Energy Management". A.Tang, S. Sethumadhavan, and S. Stolfo, Columbia University, 26th USENIX Security Symposium, 2017

- Compromised OS kernel driver used to operate circuit beyond DVFS limits
- Faults are injected which reveal cryptographic keys and perform untrusted software download
- Attack demonstrated on ARM Trustzone in a Nexus mobile phone

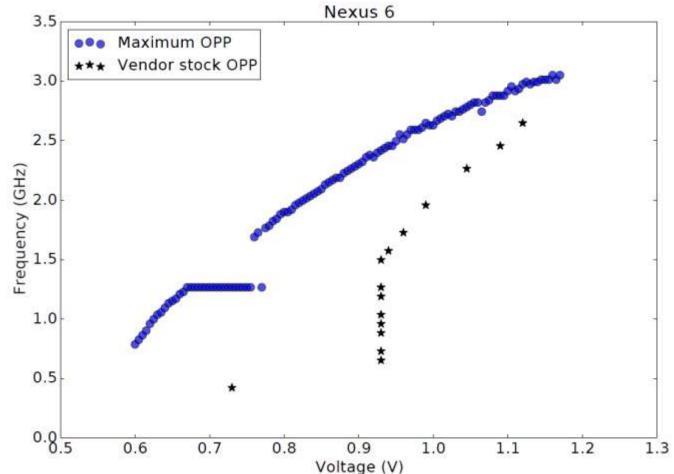
The concept of voltage and frequency scaling

- Voltage reduction saves power, but
- Voltage reduction reduces performance
- \rightarrow Operating frequency must be reduced
- If operating frequency is too high for a given voltage, faults occur
- Registers configure voltage and frequency
- In the attacked implementation voltage is configured for the complete system, but
- Clocks can be configured for each core individually



Example attacks performed with CLKskrew

- ARM TrustZone is supposed to provide an isolated trusted execution environment
- The OS kernel is able to change the clock and voltage values even when TZ code is running
- Two sample attacks have been performed:
 - Attacking AES keys of decryptions executed in TrustZone
 - Signature verification of software downloads



Intermediate conclusion

The risk is increasing!

- Increasing number of
 - Software controlled hardware features
 - Hardware functions at the reliability border
- Every software controlled hardware feature is a target for remote fault and side channel attacks
- Through side channels almost everything is observable